

# SOLUTION TO ASSIGNMENT 1

2012/2013 Semester I

MA4264

Game Theory

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1. Question 4 part 2, in Tutorial set 1;
2. Question 3 part 2, in Tutorial set 2;
3. Question 8, in Tutorial set 2.

**Solution for Question 1** Under the new payoff rules, the best response becomes:

$$R_i(s_j) = \begin{cases} \emptyset, & \text{if } s_j < \frac{1}{\sqrt{2}}; \\ [0, 1], & \text{if } s_j \geq \frac{1}{\sqrt{2}}, \end{cases}$$

where  $(i, j) = (1, 2)$  or  $(2, 1)$ . Note that when  $s_j < \frac{1}{\sqrt{2}}$ , Player  $i$  does not have the best response, because he will try to choose  $s_i$  as close as possible to  $\sqrt{1/2 - s_j^2}$ , but can not achieve  $\sqrt{1/2 - s_j^2}$ . The detailed discussion is as follows:

- For any  $1 \geq s_i \geq \sqrt{1/2 - s_j^2}$ , Player  $i$ 's payoff is 0, which is less than the payoff when Player  $i$  chooses  $\frac{1}{2}\sqrt{1/2 - s_j^2}$ ; Hence such a  $s_i$  can not be a best response.
- For any  $0 \leq s_i < \sqrt{1/2 - s_j^2}$ , Player  $i$ 's payoff is  $s_i$ , which is less than the payoff when Player  $i$  chooses  $\frac{s_i + \sqrt{1/2 - s_j^2}}{2}$ ; Hence such a  $s_i$  can not be a best response.

Therefore, the pure-strategy Nash equilibria are

$$\left[ \frac{1}{\sqrt{2}}, 1 \right] \times \left[ \frac{1}{\sqrt{2}}, 1 \right].$$

**Solution for Question 2** If  $0 < c_1 < c_2 < a$  and  $2c_2 > a + c_1$ , then  $a - c_1 > a - c_2 > \frac{a - c_2}{2} > 0$  and  $\frac{a - c_1}{2} > a - c_2 > 0$ . Hence we have the Figure (1), and from it we will obtain the Nash equilibrium:  $(\frac{a - c_1}{2}, 0)$ .

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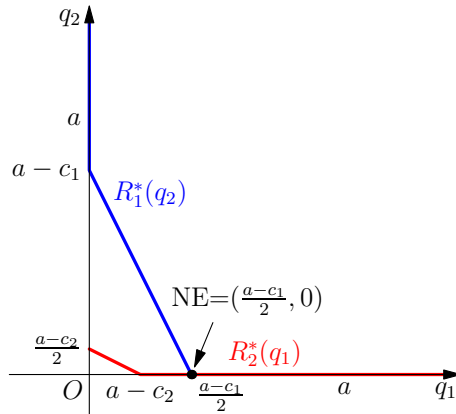


Figure 1: Intersection of best-response correspondences

**Solution for Question 3**

- (i) (1)  $C$  is strictly dominated by  $A$  and will be eliminated;
- (2)  $L$  is strictly dominated by  $M$  and will be eliminated;
- (3)  $B$  is strictly dominated by  $D$  and will be eliminated.

Hence we will obtain the reduced game  $G_1$ .

		Player 2	
		$M$	$R$
Player 1	$A$	2, 5	2, 0
	$D$	1, 1	3, 3
		$G_1$	

		Player 2	
		$M$	$R$
Player 1	$A$	2, 5	2, 0
	$D$	1, 1	3, 3
		$G_2$	

- (ii) From the bi-matrix  $G_2$ , we obtain the pure-strategy Nash equilibria:  $(A, M)$  and  $(D, R)$  (red pairs) with payoffs  $(2, 5)$  and  $(3, 3)$ , respectively.
- (iii) Let  $p_1 = (r, 1 - r)$  be a mixed strategy in which Player 1 plays  $A$  with probability  $r$ . Let  $p_2 = (q, 1 - q)$  be a mixed strategy in which Player 2 plays  $M$  with probability  $q$ . Then Player 1's expected payoff is:

$$U_1(A, p_2) = 2q + 2(1 - q) = 2,$$

$$U_1(D, p_2) = q + 3(1 - q) = 3 - 2q.$$

Hence

$$r^*(q) \equiv \arg \max_{0 \leq r \leq 1} U_1(p_1, p_2) = \begin{cases} \{1\}, & \text{if } q > \frac{1}{2}; \\ \{0\}, & \text{if } q < \frac{1}{2}; \\ [0, 1], & \text{if } q = \frac{1}{2}. \end{cases}$$

Similarly, Player 2's expected payoff is:

$$U_2(p_1, M) = 5r + (1 - r) = 1 + 4r,$$

$$U_2(p_1, R) = 3(1 - r).$$

Hence

$$q^*(r) \equiv \arg \max_{0 \leq q \leq 1} U_2(p_1, p_2) = \begin{cases} \{1\}, & \text{if } r > \frac{2}{7}; \\ \{0\}, & \text{if } r < \frac{2}{7}; \\ [0, 1], & \text{if } r = \frac{2}{7}. \end{cases}$$

We draw the graphs of  $r^*(q)$  and  $q^*(r)$  together:

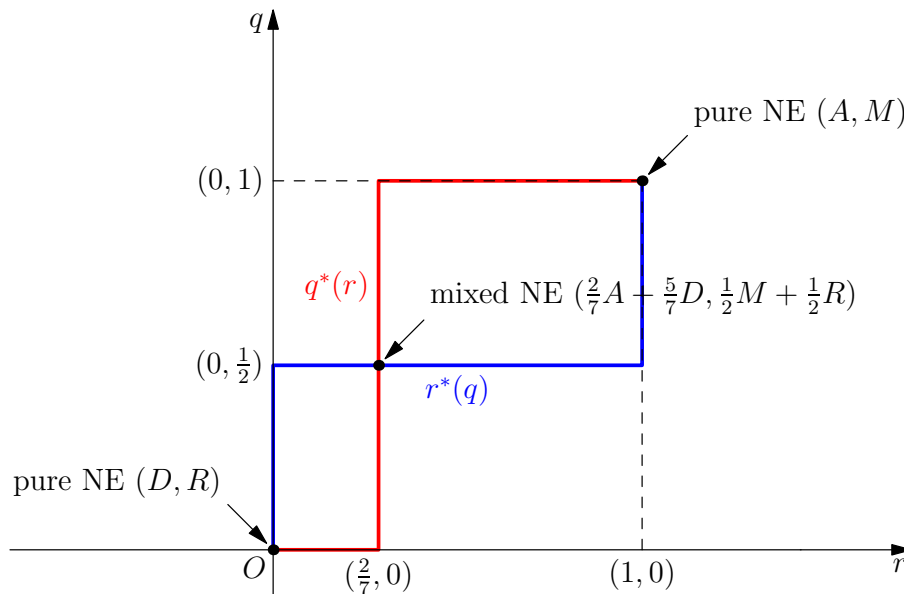


Figure 2: Intersection of best-response correspondences

The graphs of the best response correspondences  $r^*(q)$  and  $q^*(r)$  intersect at 3 points ( $r = \frac{2}{7}, q = \frac{1}{2}$ ),  $(0, 0)$  and  $(1, 1)$ . Hence, there are 3 mixed-strategy Nash equilibria:

- $(1A, 1M)$  (or  $((1, 0, 0, 0), (0, 1, 0, 0))$ ) with expected payoff  $(2, 5)$ ,
- $(1D, 1R)$  (or  $((0, 0, 0, 1), (0, 0, 1, 0))$ ) with expected payoff  $(3, 3)$ ,
- $(\frac{2}{7}A + \frac{5}{7}D, \frac{1}{2}M + \frac{1}{2}R)$  (or  $((\frac{2}{7}, 0, 0, \frac{5}{7}), (0, \frac{1}{2}, \frac{1}{2}, 0))$ ), with expected payoff  $(2, \frac{15}{7})$ .